



## A New Model for Measuring the Performance Cost of Deadline Misses

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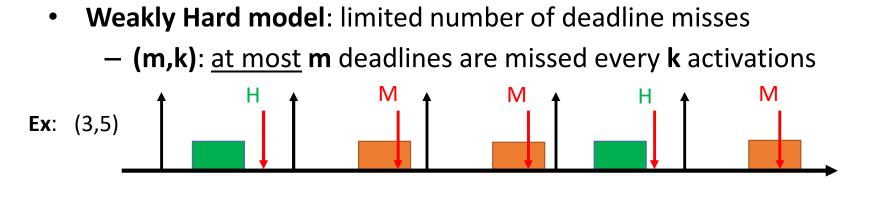
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### Introduction

- Embedded systems with control tasks may face **overload** conditions (e.g. automotive)
- <u>Common (practical) approach</u>: running at a high rate and allowing some **deadline miss** is an acceptable compromise
- Missing (few) deadlines: not catastrophic!

#### How to study performance evolution under overload conditions?





### Weakly hard model limitations

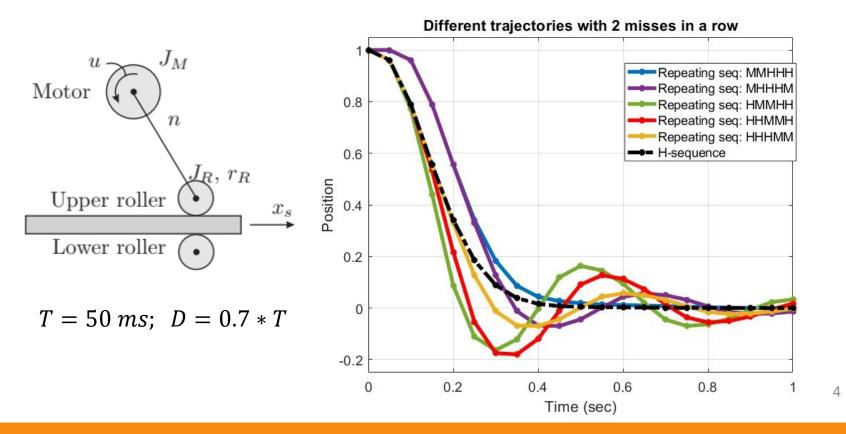
- (m,k) constraint is not enough descriptive...
- (m,k) constraint leads to a **binary** model (either pass or fail)
  - Easy to define stability guarantees
  - No information about performance of different patterns
  - Difficult to extract an **ordering** between constraints
- No relation with the system state:
  - Deadline misses may have different effects (transients vs steady state)



### Weakly hard model limitations

Different patterns of H/M deadlines lead to different performance evolutions!

Assumption: When a deadline is missed, the control output is not updated





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### A new model for performance analysis

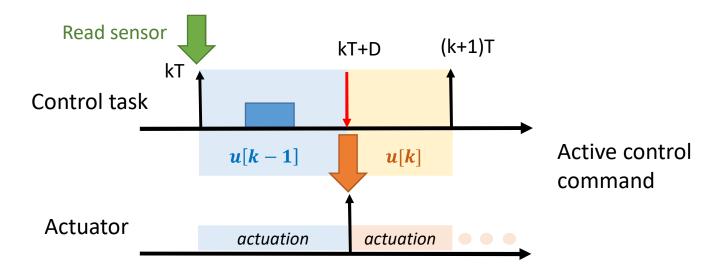
- <u>Goal:</u> Developing a new model for studying:
  - How the performance changes with different patterns of missed deadlines that satisfy a given (m,k) constraint
  - Worst guaranteed performance
  - Different policy at deadline miss (continue or kill?)
- Merging real-time analysis with control system dynamics and performance analysis





### **System model**

- Linear Time Invariant plant, MIMO
- Periodic control of period  $T_i$  and deadline  $D_i \leq T_i$
- State-feedback control: u[k] = K(r[k] x[k])



State update function:  $x[k + 1] = A_d x[k] + B_{d1}u[k - 1] + B_{d2}u[k]$ 

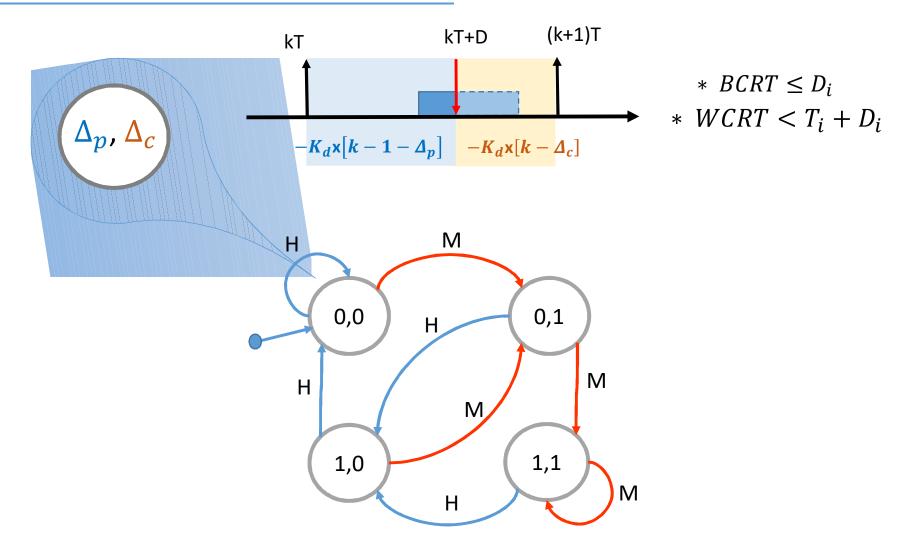
• Similar to LET model: trading jitter for latency

# **Missing a deadline** kT kT+D (k+1)TControl task u[k-1] u[k]?

- Missing a deadline means missing an actuator command update:
  <u>Keep the previous actuation value</u>
- The system dynamics changes!

• Update freshness  $\Delta$  (ageing steps) of the control output

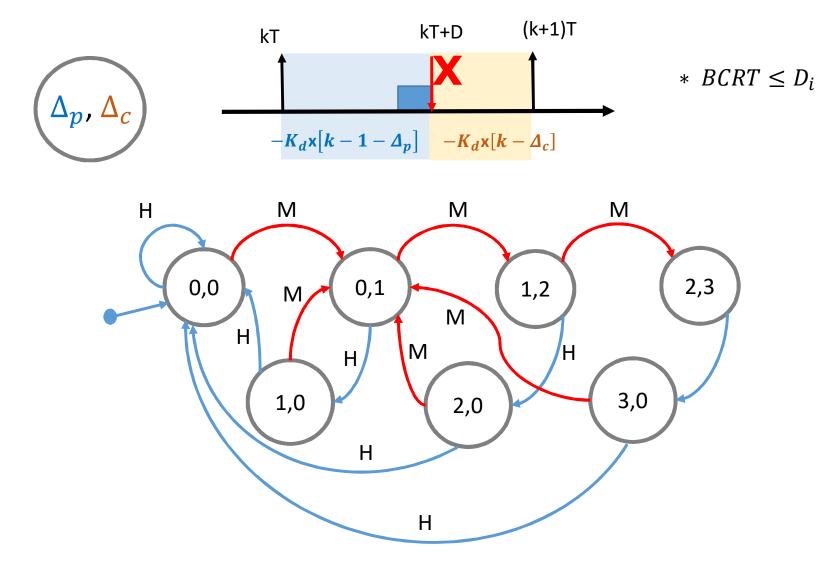
### **Update freshness: Continue strategy**





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### **Update freshness: Kill strategy**



In this example, maximum number of consecutive deadline misses is equal to 3



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### **State update matrix**

- System dynamics as a function of freshness pairs  $x[k+1] = A_d x[k] - B_{d1} K_d x[k-1-\Delta_p] - B_{d2} K_d x[k-\Delta_c]$
- Augmented state vector  $\xi[k]$

$$\begin{split} \boldsymbol{\xi}[k] &= [\boldsymbol{x}[k]; \boldsymbol{x}[k-1]; \dots \boldsymbol{x}[k-\Delta_{max}-1]] \\ & \quad \boldsymbol{\xi}[k+1] = \boldsymbol{\Phi}(\boldsymbol{\varDelta}_p, \boldsymbol{\varDelta}_c \;) \; \boldsymbol{\xi}[k] \end{split}$$

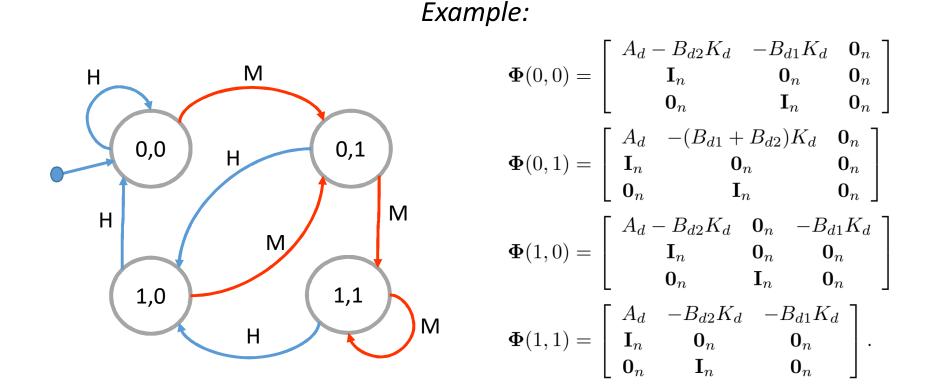
• State update matrix  $\Phi(\Delta_p, \Delta_c)$ 

$$\Phi(\Delta_{p}, \Delta_{c}) = \begin{bmatrix} A_{d} & \cdots & -B_{d2}K_{d} & \cdots & -B_{d1}K_{d} & \cdots \\ I_{n} & 0_{n} & \cdots & \cdots & \cdots \\ 0_{n} & I_{n} & 0_{n} & \cdots & \cdots \\ \vdots & \vdots & \vdots & \ddots & \cdots & \cdots \end{bmatrix}$$



### **State update matrix: an example**

• Every combination of  $(\Delta_p, \Delta_c)$  is mapped to a specific dynamic of the system through the matrix  $\Phi(\Delta_p, \Delta_c)$ 

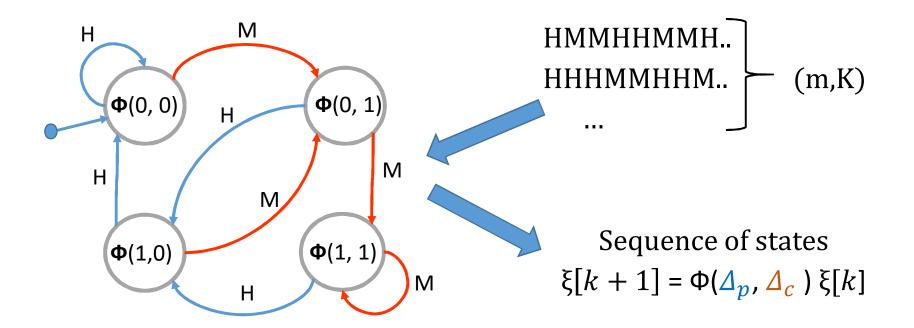


#### **Constrained switched linear system**



### **State update matrix: an example**

• Every combination of  $(\Delta_p, \Delta_c)$  is mapped to a specific dynamic of the system through the matrix  $\Phi(\Delta_p, \Delta_c)$ *Example:* 



**Constrained switched linear system** 



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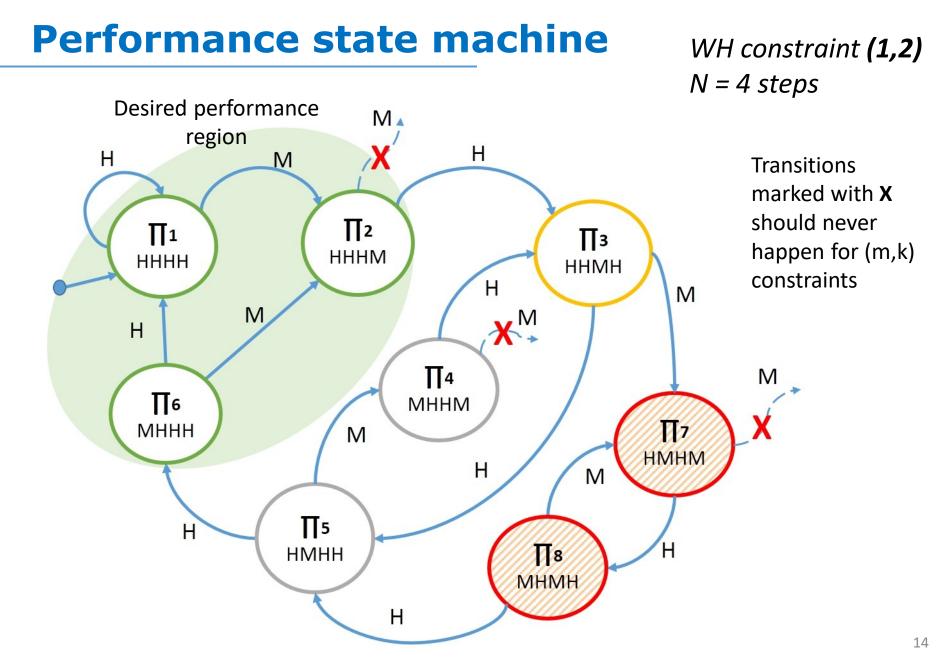
### **Performance analysis**

- Assign a **performance value** for each sequence of N jobs
- Sum of quadratic error

$$P(s) = \sum_{i=0}^{N-1} \xi[i]^T \xi[i]$$
  
=  $\xi[0]^T \left( \mathbf{I} + \mathbf{\Phi}_0^T \mathbf{\Phi}_0 + \mathbf{\Phi}_0^T \mathbf{\Phi}_1^T \mathbf{\Phi}_1 \mathbf{\Phi}_0 + \dots + \mathbf{\Phi}_0^T \mathbf{\Phi}_1^T \cdots \mathbf{\Phi}_{N-1}^T \mathbf{\Phi}_{N-1} \cdots \mathbf{\Phi}_1 \mathbf{\Phi}_0 \right) \xi[0]$   
=  $\xi[0]^T \mathbf{\Psi}(s) \xi[0]$ 

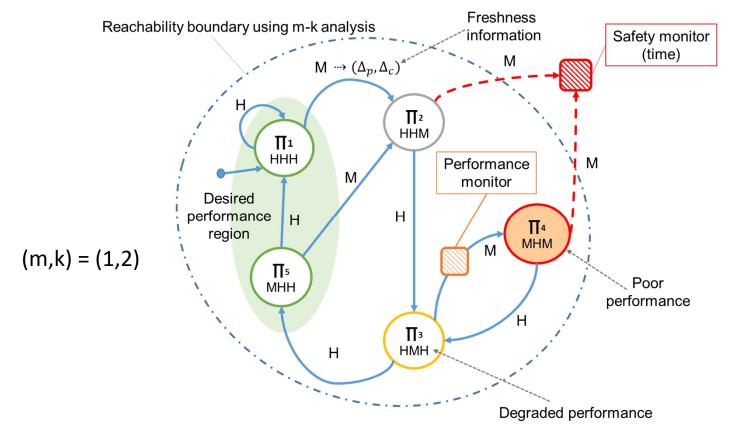
- Matrix elements of  $\Psi(s)$  depends on the **ordered** sequence of H/M
- $\prod(s) = ||\Psi(s)||_2$
- Worst Case Normalized Performance:  $WCPn = \frac{max_s \prod(s)}{\prod(all \ hits)}$





### **Possible applications**

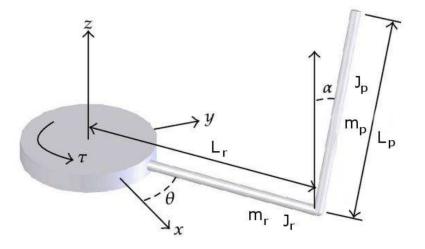
- This new model can be used as a **time contract** between software designers and control engineers
- Possibility of inserting run-time monitors





### **Case study: Furuta pendulum**

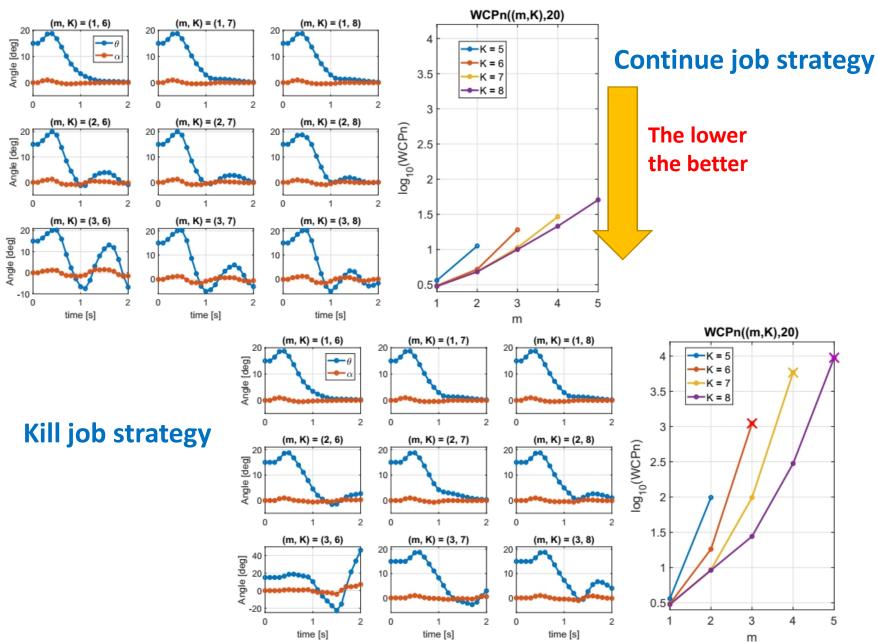
• Furuta pendulum: rotary inverted pendulum



- Linearized model in the neighbourhood of the upward position
- Feedback control with  $T_i = 0.1 sec$  and  $D_i = 0.2 * T_i$
- Testing different (m,K) values and studying how Worst Case performance changes







### **Case study: Furuta pendulum**

### **Summary**

- New model for studying performance evolution under overload conditions
- 1. Creating a state machine for computing **freshness** of outputs, applicable to different patterns and handling of deadline misses
- 2. Intergrating freshness information with state evolution of the controlled system: different **operating modes**
- 3. Creating a state machine for computing **performance** values realted to patterns of H/M deadlines
  - Worst case performance guarantees
  - Runtime monitors for performance evolution
- Case study: Furuta pendulum

### **Future work**

- Extensions:
  - Including additional performance metrics
  - Extending the case study to WCRT>T+D, allowing multiple pending jobs at deadline
- Finding **optimal controller** for a system under (m,K) constraints, for achieveing a given performance
- Adaptive control when deadline misses occur
- More complex case studies:
  - Testing non linear systems performance by simulation
  - More complex deadline miss handlings





# Thank you!

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More details in

Paolo Pazzaglia, Luigi Pannocchi, Alessandro Biondi, and Marco Di Natale,

"Beyond the Weakly Hard Model: Measuring the Performance Cost of Deadline Misses",

Proceedings of the 30th Euromicro Conference on Real-Time Systems (ECRTS 18), Barcelona, Spain, July 3-6, 2018.



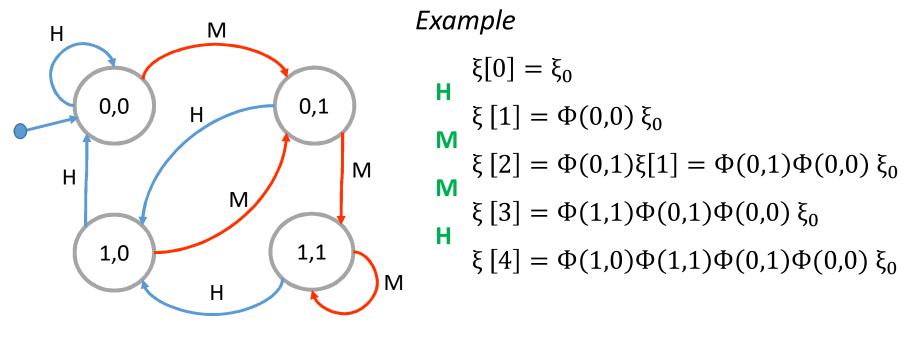
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### **Performance analysis**

- How performance changes for different patterns of H/M deadlines?
- H/M patterns are mapped to state trajectories of the system

State update equation:  $\xi[k+1] = \Phi(\Delta_p, \Delta_c) \xi[k]$ 



General state trajectory equation:  $\xi[k+1] = \Phi_k \Phi_{k-1} \dots \Phi_0 \xi_0$ 

